

Lecture 14 - Nov 3

Graphs

Tracing BFS using a FIFO Queue

Back Edge (DFS) vs. Cross Edge (BFS)

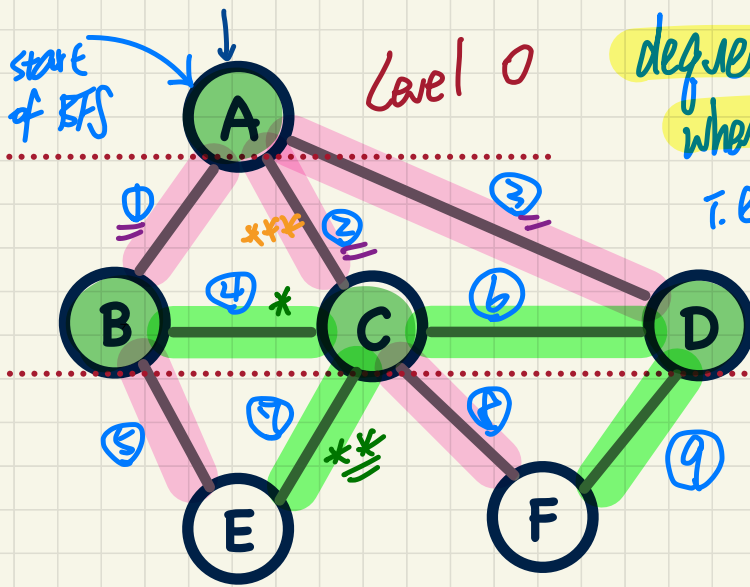
Implementing Graphs: Adjacency Lists

Announcements/Reminders

- Today's class: notes template posted
- **Test 1** results to be released on Tuesday (Nov 4)
- Change of Dates:
 - + **Assignment 2** to be released on Wed, Nov 12
 - + **Assignment 2** to be due on Wed, Nov 19
 - + **Test 2** to be take place on Mon, Nov 24

*** not possible to have a cross edge linking back to a previous level

Breadth-First Search (BFS): Example 1 (a)



dequeue a vertex when all its i.e. have been marked.

enqueued
A
B
C
D
E
F

dequeued
A
B
C
D
E
F

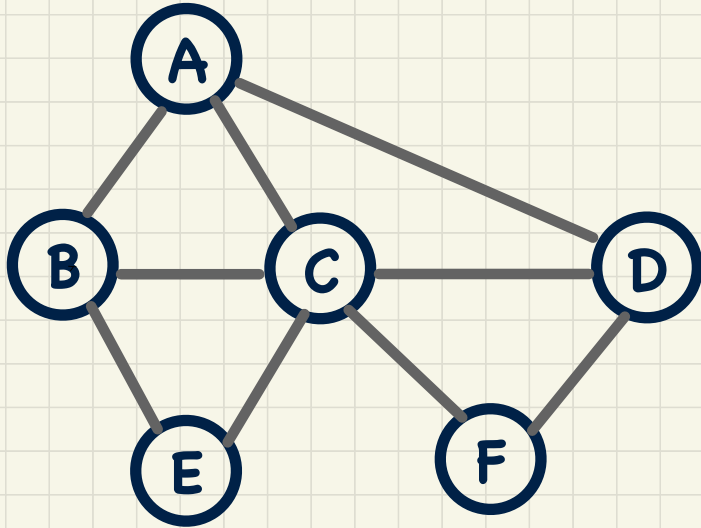
* Cross edge connecting vertices at the same level.
** cross edge connecting vertices at the next level.

Assumptions:

- Adjacent vertices visited in alphabetic order



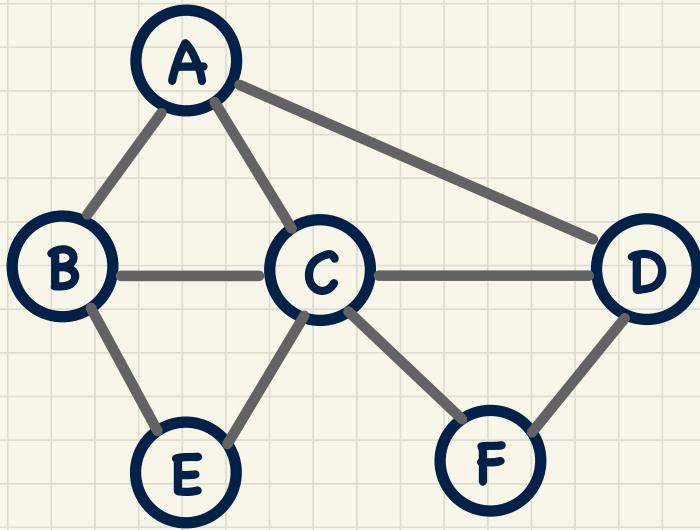
Breadth-First Search (BFS): Example 1 (b)



Assumptions:

- Adjacent vertices visited in alphabetic order
- Exception: Edge AC visited first

Breadth-First Search (BFS): Example 1 (c)

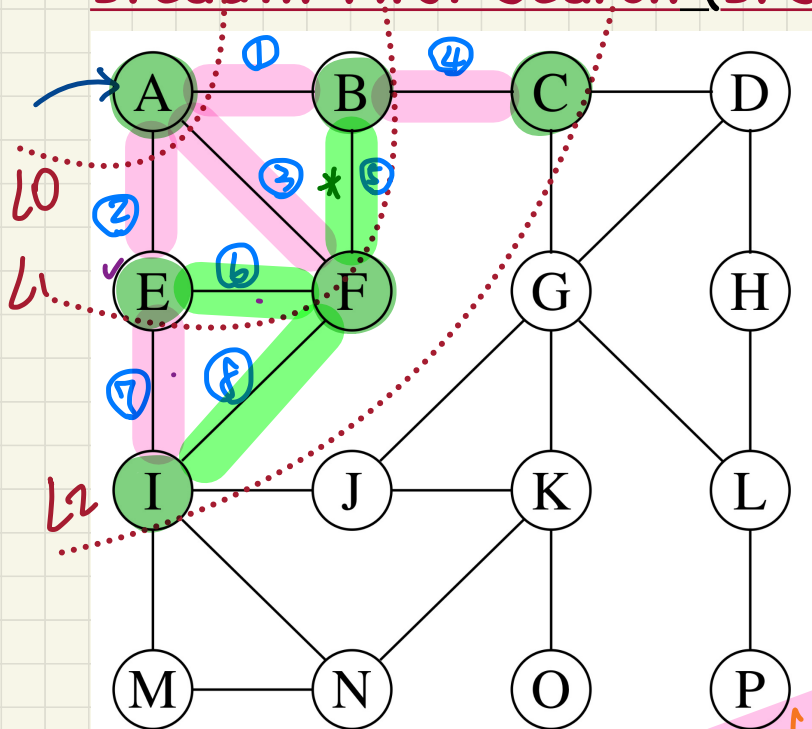


Assumptions:

- Adjacent vertices visited in alphabetic order
- Exception: Edge AD visited first

Breadth-First Search (BFS): Example 2

* Though organized vertically, vertices B and I are at the same level.



enqueued	dequeued
A	A
B ①	B
E ②	E
F ③	F
C ④	

front

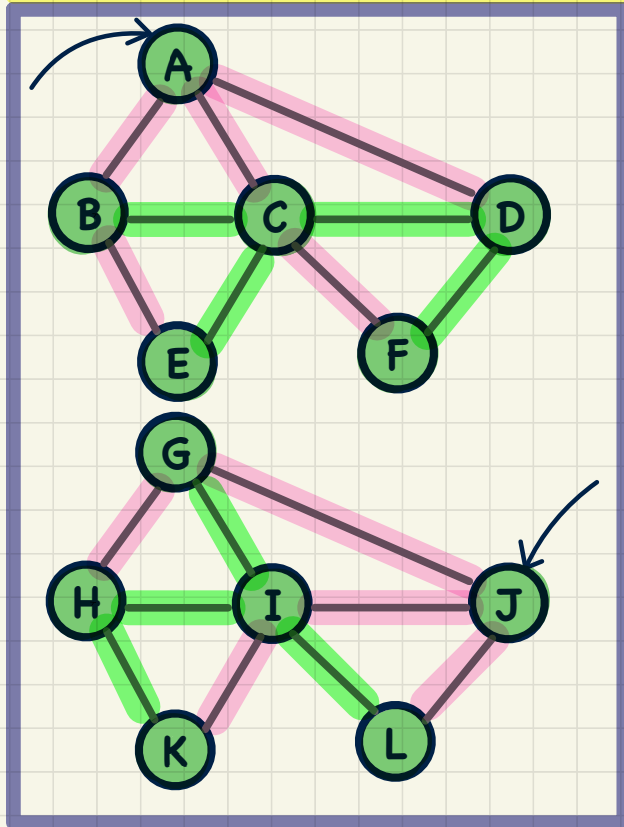
① ② ③ ④ ⑦

~~A~~ ~~B~~ ~~E~~ ~~F~~ C I

next step of BFS:
go over each vertex at L2 and reach out to L3

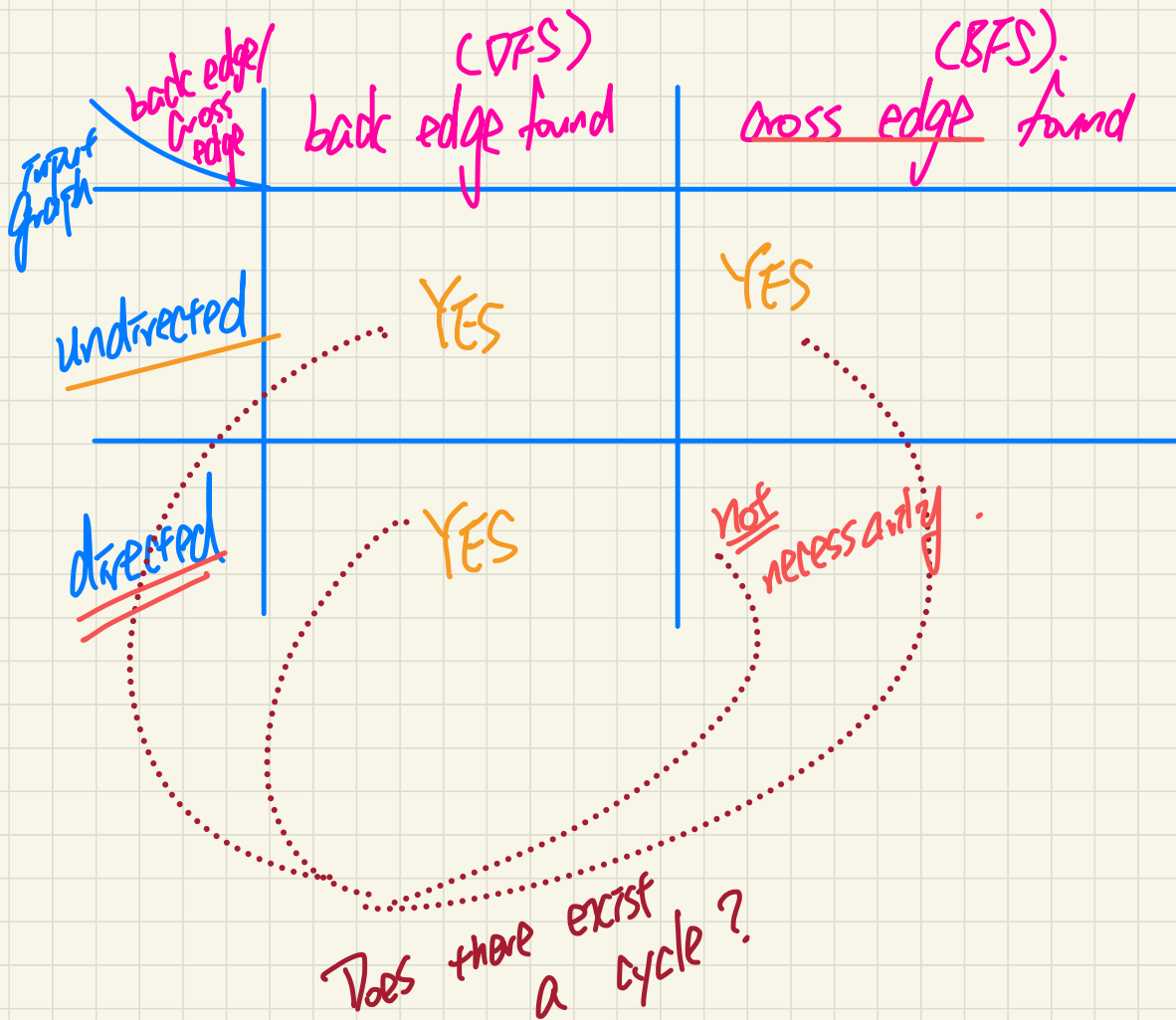
Graph Traversals: Adapting BFS

Efficient Traversal of Graph G:



Graph Questions:

- Find a **path** between $\{u, v\} \subseteq V$
Start BFS from u, maintain the path until v is encountered.
 - Is v **reachable** from u?
Start BFS from u, is v ever encountered?
 - Find **all connected components** of G.
Keep doing BFS, until all vertices are visited.
 - Compute a **spanning tree** of a connected G.
Each BFS will identify the spanning tree containing the start vertex
 - Is G **connected**?
How many BFSes are needed to cover all vertices?
 - If G is cyclic, return a **cycle**.
- ①



int numVertexVisited = 0 ;
while (numVertexVisited < |V|) { exit loop as soon as num V. Visited $\geq$ |V|

Pick some vertex π (unvisited),
start BFS on π .

↓
numVertexVisited properly updated
(for every discovery edge)

iterations
||
CCs.

}

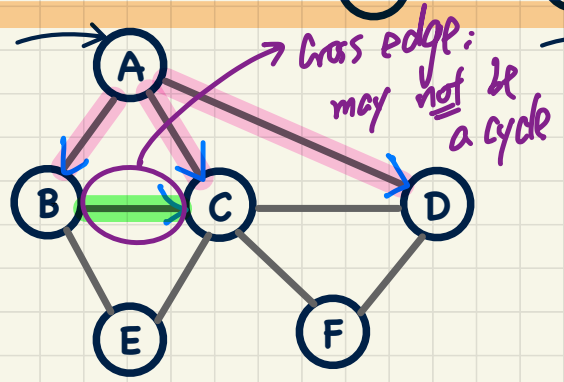
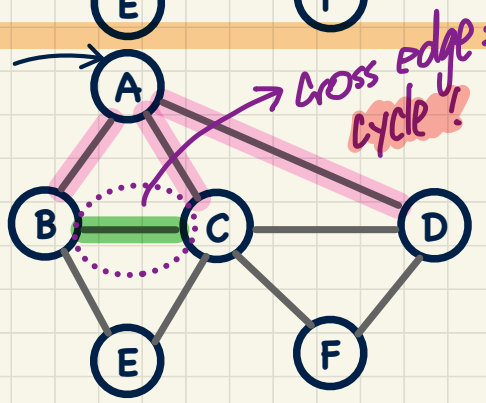
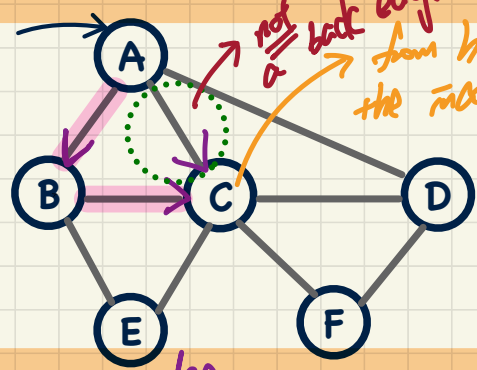
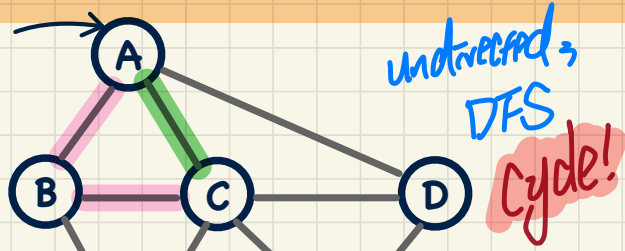
Back Edge (DFS) vs. Cross Edge (BFS): Cyclic?

focus on finding a path leading to a back edge.

YES.

Does a back edge always imply the existence of a cycle?

Does a cross edge always imply the existence of a cycle?



Graphs in Java: Adjacency List Strategy (1)

```
class AdjacencyListGraph<V, E> implements Graph<V, E> {  
    private DoublyLinkedList<AdjacencyListVertex<V>> vertices;  
    private DoublyLinkedList<AdjacencyListEdge<E, V>> edges;  
    private boolean isDirected;  
  
    /* initialize an empty graph */  
    AdjacencyListGraph(boolean isDirected) {  
        vertices = new DoublyLinkedList<>();  
        edges = new DoublyLinkedList<>();  
        this.isDirected = isDirected;  
    }  
}
```

```
public class Vertex<V> {  
    private V element;  
    public Vertex(V element) { this.element = element; }  
    /* setter and getter for element */  
}
```

```
public class Edge<E, V> {  
    private E element;  
    private Vertex<V> origin;  
    private Vertex<V> dest;  
    public Edge(E element) { this.element = element; }  
    /* setters and getters for element, origin, and destination */  
}
```

```
public class EdgeListVertex<V> extends Vertex<V> {  
    public DLNode<Vertex<V>> vertexListPosition;  
    /* setter and getter for vertexListPosition */  
}
```

```
public class EdgeListEdge<E, V> extends Edge<E, V> {  
    public DLNode<Edge<E, V>> edgeListPosition;  
    /* setter and getter for edgeListPosition */  
}
```

```
class AdjacencyListVertex<V> extends EdgeListVertex<V> {  
    private DoublyLinkedList<AdjacencyListEdge<E, V>> incidentEdges;  
    /* getter for incidentEdges */  
}
```

```
class AdjacencyListEdge<V> extends EdgeListEdge<V> {  
    DLNode<Edge<E, V>> originIncidentListPos;  
    DLNode<Edge<E, V>> destIncidentListPos;  
}
```

for efficient
removal
of edges